
Lecture 09X: C Function Pointers

Concurrent and Multicore Programming


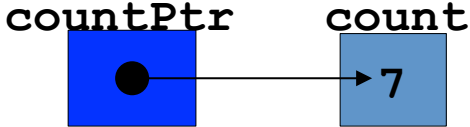
Department of Computer Science and Engineering

Yonghong Yan

yan@oakland.edu

www.secs.oakland.edu/~yan

Pointer Variable Declarations and Initialization

- Pointer variables
 - Contain memory addresses as their values
 - Normal variables contain a specific value (direct reference)
 - `int count = 7;` 
 - Pointers contain address of a variable that has a specific value (indirect reference)
 - Indirection – referencing a pointer value
 - `int count = 7;`
 - `int * countPtr = &count;` 

Pointer Variable Declarations and Initialization

- Pointer declarations
 - ***** used with pointer variables

```
int *myPtr;
```
 - Declares a pointer to an **int** (pointer of type **int ***)
 - Multiple pointers require using a ***** before each variable declaration

```
int *myPtr1, *myPtr2;
```
 - Can declare pointers to any data type
 - Initialize pointers to **0**, **NULL**, or an address
 - **0** or **NULL** – points to nothing (**NULL** preferred)

Pointer Operators

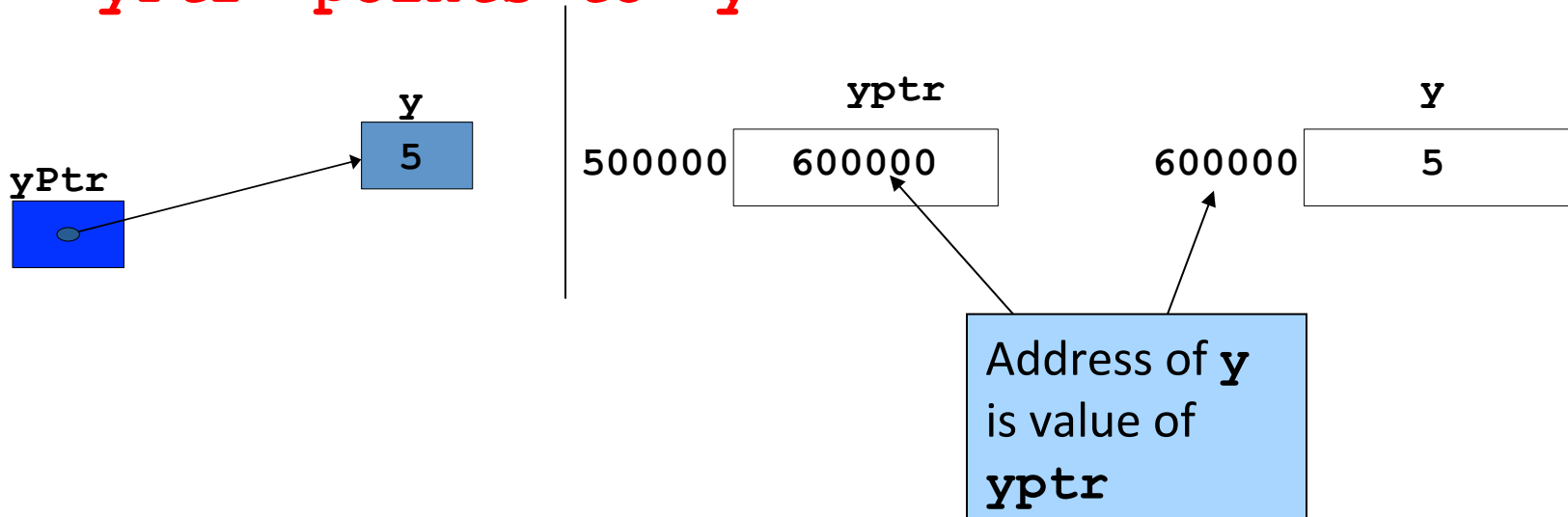
- `&` (address operator)
 - Returns address of operand

```
int y = 5;
```

```
int *yPtr;
```

```
yPtr = &y;           // yPtr gets address of y
```

```
yPtr "points to" y
```



Pointers and Arrays

- Arrays and pointers closely related
 - Array name like a constant pointer
 - Pointers can do array subscripting operations
- Declare an array **b [5]** and a pointer **bPtr**
 - To set them equal to one another use:
 - `bPtr = b;`
 - The array name (**b**) is actually the address of first element of the array **b [5]**
 - `bPtr = &b [0]`
 - Explicitly assigns **bPtr** to address of first element of **b**

Pointers and Arrays

- Element **b [3]**
 - Can be accessed by *** (bPtr + 3)**
 - Where **n** is the offset. Called pointer/offset notation
 - Can be accessed by **bptr [3]**
 - Called pointer/subscript notation
 - **bPtr [3]** same as **b [3]**
 - Can be accessed by performing pointer arithmetic on the array itself
 - * (b + 3)**

Pointers to Functions

- Pointer to function
 - Contains address of function
 - Similar to how array name is address of first element
 - Function name is starting address of code that defines function
- Function pointers can be
 - Passed to functions
 - Stored in arrays
 - Assigned to other function pointers

Pointers to functions: Variable for functions

- Declaration:

```
returnType (*funVarName)(parameterTypes);
```

- Examples:

```
int (*f)(int, float);
```

pointer to a function that takes an integer argument and a float argument and returns an integer

```
int *(*g[])(int, float);
```

pointer to a function that takes an integer argument and a float argument and returns a *pointer* to an integer

```
int *(*g[])(int, float);
```

An *array* of pointers to functions – Each function takes an integer argument and a float argument and returns a pointer to an integer

Pointers to functions: WHY?

- They allow for a certain amount of **polymorphism**:
 - “poly” (many) + “morph” (shape)
 - A polymorphic language can handle a range of different data types (“shapes”?) with a single statement
- This is common in OO languages like C++, Java:

```
Animal myPet;  
...  
myPet.makeSound();
```

This method call will result in different sounds, depending on whether `myPet` holds a `COW` object, an `Elephant` object, etc.

Example: searching a singly-linked list

```
typedef struct IntNode {  
    int value;          struct IntNode *next;  
} INTNODE;
```

OK, but it only works for nodes containing integer data. If you want a list of strings, you'll need to define a new type and new function.

```
INTNODE *search_list(INTNODE *node, int const key) {  
    while (!node) {  
        if (node->value == key) break;  
        node = node->next;  
    }  
    return node;  
}
```

A more abstract notion of “node”

```
typedef struct Node {  
    void *value; struct Node *next;  
} NODE;
```

`void*` is compatible with any pointer type.
So, this member can hold (a pointer to) any value!

```
void construct_node(NODE *node, void *value, NODE *next) {  
    node->value = value; node->next = next;  
}
```

```
NODE *new_node(void *value, NODE *next) {  
    NODE *node = (NODE *)malloc(sizeof(NODE));  
    construct_node(node, value, next);  
    return node;  
}
```

A more abstract notion of “search list”

- What is it that makes the old `search_list` only work for integers?
 - The key parameter is of type `int`
 - The `==` operator is used to compare `int` values – but `==` will not work for many types (e.g. structs, strings)
- A solution: pass in an additional argument – a comparison function!
 - Programmer must supply a comparison function that’s appropriate for the data type being stored in the nodes
 - This function argument is called a **callback function**:
 - Caller passes in a pointer to a function
 - Callee then “calls back” to the caller-supplied function

Abstract “search list” with callback function

```
NODE *search_list(NODE *node, void const *key,  
    int (*compare)(void const *, void const *)) {  
  
    while (node) {  
        if (!compare(node->value, key)) break;  
        node = node->next;  
    }  
    return node;  
  
}
```

Assumption: compare returns zero if its parameter values are equal; nonzero otherwise

Using callback functions

- If our nodes hold strings, we have a compare function already defined: `strcmp` or `strncmp`

```
#include <string.h>
```

```
...
```

```
match = search_list(root, "key", &strcmp);
```

& is optional here –
compiler will implicitly take the address

Note: you may get a warning, since `strcmp` is not strictly of the right type:
its parameters are of type `char *` rather than `void *`

Using callback functions

- If our nodes hold other kinds of data, we may need to “roll our own” compare function

```
int compare_ints(void const *a, void const *b) {  
    const int ia = *(int *)a, ib = *(int *)b;  
    return ia != ib;  
}
```

...

```
match = search_list(root, key, &compare_ints);
```

Jump tables

- In some cases, a nice alternative to long, repetitive switch statements, like this:

```
double add(double, double);  
double sub(double, double);  
double mul(double, double);  
double div(double, double);
```

```
switch(oper) {  
case ADD:    result = add(op1, op2); break;  
case SUB:    result = sub(op1, op2); break;  
case MUL:    result = mul(op1, op2); break;  
case DIV:    result = div(op1, op2); break;  
}
```


Jump tables

- Jump table alternative:

```
double add(double, double);  
double sub(double, double);  
double mul(double, double);  
double div(double, double);
```

Array of pointers to functions.
Each function takes two `double`s
and returns a `double`

```
double (*oper_func[])(double, double) = {  
    add, sub, mul, div  
};
```

```
result = oper_func[oper](op1, op2);
```

Pointers to functions: safety concerns

- What if uninitialized function pointer value is accessed?
 - Safest outcome: memory error, and program is terminated
 - But what if the “garbage” value is a valid address?
 - Worst case: address contains program instruction – execution continues, with random results
 - Hard to trace the cause of the erroneous behavior

References

- The Function Pointer Tutorials.
<http://www.newty.de/fpt/index.html>